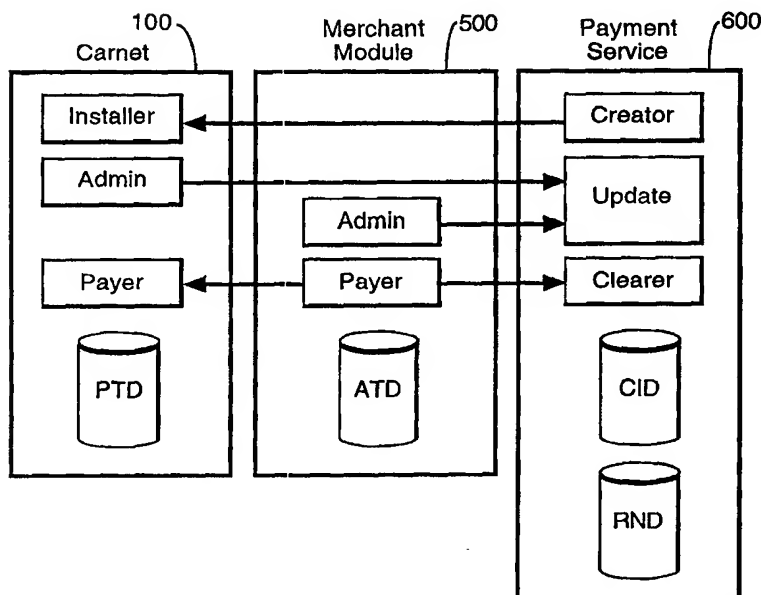




INTERNATIONAL APPLICATION PUBLISHED UNDER THE PATENT COOPERATION TREATY (PCT)

(51) International Patent Classification ⁶ : G07F 7/08, 7/10	A1	(11) International Publication Number: WO 98/22915 (43) International Publication Date: 28 May 1998 (28.05.98)
(21) International Application Number: PCT/GB97/03116 (22) International Filing Date: 12 November 1997 (12.11.97) (30) Priority Data: 9624127.8 20 November 1996 (20.11.96) GB (71) Applicant (for all designated States except US): BRITISH TELECOMMUNICATIONS PUBLIC LIMITED COMPANY [GB/GB]; 81 Newgate Street, London EC1A 7AJ (GB). (72) Inventor; and (75) Inventor/Applicant (for US only): HILL, Jake [GB/GB]; Parkside, Hackney Road, Woodbridge, Suffolk IP12 1NW (GB). (74) Agent: WELLS, David; BT Group Legal Services, Intellectual Property Dept., 8th floor, Holborn Centre, 120 Holborn, London EC1N 2TE (GB).		(81) Designated States: AL, AM, AT, AU, AZ, BA, BB, BG, BR, BY, CA, CH, CN, CU, CZ, DE, DK, EE, ES, FI, GB, GE, GH, HU, ID, IL, IS, JP, KE, KG, KP, KR, KZ, LC, LK, LR, LS, LT, LU, LV, MD, MG, MK, MN, MW, MX, NO, NZ, PL, PT, RO, RU, SD, SE, SG, SI, SK, SL, TJ, TM, TR, TT, UA, UG, US, UZ, VN, YU, ZW, ARIPO patent (GH, KE, LS, MW, SD, SZ, UG, ZW), Eurasian patent (AM, AZ, BY, KG, KZ, MD, RU, TJ, TM), European patent (AT, BE, CH, DE, DK, ES, FI, FR, GB, GR, IE, IT, LU, MC, NL, PT, SE), OAPI patent (BF, BJ, CF, CG, CI, CM, GA, GN, ML, MR, NE, SN, TD, TG). Published <i>With international search report.</i>

(54) Title: TRANSACTION SYSTEM



(57) Abstract

In a digital payment system a sequence of random numbers are stored at a payment server (600). A set of digitally encoded random numbers derived from the stored sequence are issued to the user in return for payment. The tokens are stored in a Carnet (100). The user can then spend the tokens by transferring tokens to a merchant (500), for example to an on-line service provider. The merchant returns each token received to the payment server. The payment server authenticates the token and transmits an authentication message to the merchant. The merchant, payment server and user may be linked by Internet connections.

FOR THE PURPOSES OF INFORMATION ONLY

Codes used to identify States party to the PCT on the front pages of pamphlets publishing international applications under the PCT.

AL	Albania	ES	Spain	LS	Lesotho	SI	Slovenia
AM	Armenia	FI	Finland	LT	Lithuania	SK	Slovakia
AT	Austria	FR	France	LU	Luxembourg	SN	Senegal
AU	Australia	GA	Gabon	LV	Latvia	SZ	Swaziland
AZ	Azerbaijan	GB	United Kingdom	MC	Monaco	TD	Chad
BA	Bosnia and Herzegovina	GE	Georgia	MD	Republic of Moldova	TG	Togo
BB	Barbados	GH	Ghana	MG	Madagascar	TJ	Tajikistan
BE	Belgium	GN	Guinea	MK	The former Yugoslav Republic of Macedonia	TM	Turkmenistan
BF	Burkina Faso	GR	Greece			TR	Turkey
BG	Bulgaria	HU	Hungary	ML	Mali	TT	Trinidad and Tobago
BJ	Benin	IE	Ireland	MN	Mongolia	UA	Ukraine
BR	Brazil	IL	Israel	MR	Mauritania	UG	Uganda
BY	Belarus	IS	Iceland	MW	Malawi	US	United States of America
CA	Canada	IT	Italy	MX	Mexico	UZ	Uzbekistan
CF	Central African Republic	JP	Japan	NE	Niger	VN	Viet Nam
CG	Congo	KE	Kenya	NL	Netherlands	YU	Yugoslavia
CH	Switzerland	KG	Kyrgyzstan	NO	Norway	ZW	Zimbabwe
CI	Côte d'Ivoire	KP	Democratic People's Republic of Korea	NZ	New Zealand		
CM	Cameroon			PL	Poland		
CN	China	KR	Republic of Korea	PT	Portugal		
CU	Cuba	KZ	Kazakistan	RO	Romania		
CZ	Czech Republic	LC	Saint Lucia	RU	Russian Federation		
DE	Germany	LI	Liechtenstein	SD	Sudan		
DK	Denmark	LK	Sri Lanka	SE	Sweden		
EE	Estonia	LR	Liberia	SG	Singapore		

TRANSACTION SYSTEM

The present invention relates to a digital payment transaction system.

With the growth and commercialisation of the internet, there has been an
5 increasing need for technologies to allow payments to be made on-line. For
transactions of relatively high financial value this need is adequately met, for
example, by systems using electronic cheques issued by a trusted party such as a
bank. Such electronic cheques are typically validated by a signature which is
encrypted using a public key algorithm such as RSA. There is however a
10 significant computational overhead associated with the use of such algorithms.
Therefore, just as in real life a cheque is unlikely to be acceptable for a purchase of
small value because of the associated transaction costs, so also in electronic
commerce, electronic cheques are not suitable for payments of low value.

A number of proposals have been made for alternative transaction systems
15 suitable for making the so-called "micropayments" required by low-value
transactions. However, it has proved technically difficult to provide the low
processing overheads required for any micropayment technology whilst maintaining
an adequate level of security.

One example of a previous proposal for a micropayment system is that
20 developed by the US corporation Digital and known as "Millicent". This system is
described by its proponents as a lightweight protocol suitable for supporting
purchases costing less than a cent. It is based on decentralised validation of
electronic cash at the vendor's server. The digital payment tokens in this system
are termed "scrip". They are issued by a central payment service in return for
25 prepayment using a conventional payment method such as a credit card. The
vendor may then accept scrip from the user in payment for goods or more typically
for services. The vendor generates fresh scrip and returns it to the user as change
for the transaction. The scrip is authenticated and fresh scrip generated by the
vendor using a hash function. This represents a potential security weakness, in
30 that if the hash function is cracked, then the scrip would be open to forgery and
duplication. Moreover, there is a processing overhead associated with the use of
the hash function by the vendor. Although this is less than the overhead
associated, for example, with the use of a PGP-encrypted signature, it is

nonetheless a significant limitation. It is admitted by the proponents of the Millicent system that, as a result of its limited efficiency, there is a practical lower bound to the transaction values it can handle. It is suggested that this lower bound is around 1/10 of a cent. Millicent is therefore not suitable for use, for
5 example, as a charging mechanism for internet usage. The costs of packet transmission on the internet have been estimated at around 1/600th of a cent.

European patent application EP-A-0507669 discloses an example of another type of payment system, based on smart card technology. Here, rather than cryptographic security being relied upon, security is based on the physical
10 integrity of the card. A randomly selected sub-set of a number of token values is withheld, so that the presence of one of the withheld token values in a subsequent transaction can serve as an indication of attempted fraud. The set of tokens issued to a particular card is not statistically random but may, for example, all fall within a limited range of numerical values, and may be ordered in sequence
15 determined by their numerical values.

According to a first aspect of the present invention, there is provided a method of operating a digital payment transaction system comprising:

- a) storing at a payment server a sequence of random numbers;
- b) issuing to a user a set of digital payment tokens comprising a
20 sequence of digitally encoded random numbers derived from the said stored sequence of random numbers;
- c) transferring a payment token from the user to a merchant platform;
- d) transferring from the merchant platform to the payment server the payment token received from the user in step (c);
- 25 e) at the payment server, authenticating the token by comparing the value of the random number of the token from the merchant platform and a value derived from a corresponding position in the stored sequence of random numbers; and
- f) subsequently communicating an authentication message from the
30 payment server to the merchant platform.

The present invention provides a digital payment transaction system which offers improved efficiency and security and which is suitable for making micropayments, including very low value payments. This is achieved by issuing to

the user a "carnet" or set of digital payment tokens which comprises a set of random numbers in a determined sequence. The numbers in the sequence are completely unrelated to one another, and there is no correlation between the numerical value of a token and its position in the sequence. This is in contrast
5 with prior art proposals in which the tokens are related to one another by cryptographic transforms. The tokens are validated by passing them to the payment server where the token is compared with the corresponding number in the random sequence stored at the server. In this way, the system offers a high level of cryptographic security, while completely removing the processing overhead from
10 the vendor.

In the description and claims the term "random" is used to encompass both truly random numbers and pseudo-random numbers which match, to within a required level of accuracy, the statistical properties of a truly random sequence.

Preferably the payment server is remote from the merchant platform and
15 the merchant platform communicates over a communications network with the payment server. Typically, although not necessarily, all three of the user, the merchant and the payment server will be linked by internet connections.

Preferably the set of digital payment tokens is derived from the sequence of random numbers stored at the payment server by selecting part of the said
20 sequence and encoding the said part of the sequence with a key which is specific to the user.

Generating the carnet by encrypting part of the random sequence using a user-specific key ensures that the required length of the random number sequence, and the associated storage needed in the payment server, scales acceptably as the
25 number of users increases. At the same time, encryption with the user key further enhances the security of the method.

Preferably the part of the said sequence is encoded by a symmetrical block cipher. This is preferred as giving the highest level of security. Alternatively, in fields of use where it is more important to reduce the computational overhead
30 involved in creating the carnet, a keyed hash function might be used. It will be understood that, by contrast with the Millicent proposal discussed above, hashing is used only within the payment server to produce the token values from the stored random number sequence, and does not determine the relationship between

the different numbers in the carnet. The security of the system does not therefore rest upon the integrity of the hash function alone.

Preferably in step (d) of the method, the merchant communicates, together with each payment token, an authentication token from a sequence of
5 authentication tokens issues by the payment server.

Although typically there will be greater trust between the vendor and the payment server, there is still a need to provide security for the transactions between these two parties. In the preferred implementation of the present invention, a particular effective approach is to use the same mechanism as that
10 used to generate the payment tokens themselves. Then the payment server issues the merchant with a series of authentication tokens. These authentication tokens may be derived from the sequence of random numbers in the same manner as the payment tokens. The merchant can then return pairs of payment tokens and authentication tokens to the payment server for authentication. The payment
15 server may then automatically update a merchant account record after authenticating a validated payment token and authentication token received from the merchant platform.

Preferably the step of authenticating the digital payment token comprises:

i) attempting to authenticate the digital payment token against a value at
20 a position in the sequence of random numbers stored at the payment server; and

ii) when the token is not authenticated in step i), attempting to authenticate the digital payment token against one or more other values in the stored sequence, which other values fall within a predetermined maximum distance from the said position;

25 and in step (f) the authentication message indicates that the authentication is successful when the token is successfully authenticated in either step (i) or step (ii) .

The basic payment mechanism assumes that the user and the payment server going through their sequences of stored random numbers in step.
30 Sometimes however tokens may arrive at the payment out of sequence, for example because one merchant is slower than another in submitting tokens for clearance. This feature of the invention enables the system to function robustly in these circumstances. Instead of checking only the next in sequence stored

random number at the payment server, a sliding window is used to select a range (in terms of sequence position) of stored random numbers. A submitted token can be successfully validated against any stored value falling within the window (provided that stored value has not been previously used).

5 Preferably the method further comprises:

f) maintaining at the payment server a record of the current state of the set of digital payment tokens; and

g) when the digital payment tokens issued to the user are lost or corrupted, communicating data from the payment server to the user and thereby
10 updating the set of digital payment tokens to a state corresponding to that recorded at the payment server.

A further significant advantage offered by preferred implementations of the invention is that if the carnet is destroyed or stolen, then it can be recreated by the payment server.

15 Preferably the method further comprises:

issuing the user an identification number (PIN);

modifying, using the identification number, the numbers derived from the said stored sequence of random numbers ;

and further modifying, using the identification number, the digital payment
20 token which is transferred to the merchant in step (c).

This preferred feature further enhances the security of the system, without adding significantly to the processing overhead at the merchant platform.

Preferably in the step of modifying the digital payment token, the digitally encoded value issued by the payment server is combined with the identification
25 number using a Boolean logic operation, and the result of the said operation is output as the modified digital payment token. Preferably the said binary logic operation is XOR.

According to a second aspect of the present invention there is provided a method of operating a digital payment transaction system comprising:

30 a) issuing to a user a set of digital payment tokens comprising a sequence of digitally encoded random numbers

b) issuing to a merchant a set of authentication tokens comprising a sequence digitally encoded of random numbers;

- c) transferring a digital payment token from the user to a merchant;
- d) transferring the said digital payment token, and with the digital payment token transferring an authentication token, from the merchant to a payment server;
- 5 e) authenticating the digital payment token and the authentication token against records stored at the payment server; and
- f) returning an authentication message to the merchant.

The invention also encompasses merchant platforms, client platforms and payment servers for use in the methods of the first and second aspects of the
10 invention, and networks including such devices. In the example described below, the client platform is a personal computer running a web browser. Other examples of client platforms include smart cards and personal digital assistants (PDA's).

Systems embodying the present invention will now be described in further detail, by way of example only, with reference to the accompanying drawings, in
15 which:

Figure 1 is a schematic showing the main components of the payment system;

Figure 2 is a flow diagram showing the main execution loop of the payment service clearer module;

20 Figure 3 is a flow diagram showing the main execution loop of the merchant module;

Figure 4 is a flow diagram showing the process for verifying tokens received at the payment service;

Figure 5 is a flow diagram showing the updater module used to restore a
25 client or merchant token database;

Figure 6 is a flow diagram showing the collector module at the merchant;

Figure 7 is a flow diagram showing the payer module in the carnet;

Figure 8 is a schematic of a network embodying the present invention;

Figure 9 is a schematic of an alternative embodiment; and

30 Figure 10 illustrates the operation of the payment service using a sliding window.

As shown in Figure 8, a client terminal 1, which in this example is a personal computer, is connected via a modem 2 and the PSTN (public switched

telephone network) 3 to an Internet Service Provider (ISP) 4. Via the internet, the client terminal forms, at different times, connections to a payment server 6 (also termed herein the "payment service") and to an on-line merchant 5. The merchant 5 offers a service in return for payment. For example, the merchant may serve 5 HTTP (hypertext transfer protocol) pages of personalised news items at a low, fixed charge per page. Alternatively, the merchant may, for example, control a node 5a which provides access to a high-speed internet connection. Access to the node is made available to the user at a charge which may be calculated, for example, on the basis of the length of time for which the user is connected.

10 Payment tokens may be requested from the user at regular intervals, or a payment token may be collected from each packet which passes through the node. The payment system embodying the present invention is termed by the inventor "QuickPay".

In use a module which is termed the "carnet" module is installed on the 15 client terminal. This module, which is described in further detail below, includes programs which support interactions with the merchant and the payment service. In addition the terminal stores data relating to any payment tokens which are currently held by the user.

Initially, the user establishes an internet connection with the payment 20 service, and purchases tokens to a certain value. This transaction may be carried out, for example, by transmitting from the client to the payment service a request for tokens to a certain value, say £10, together with a credit card number. This number may be encrypted using any one of a number of public key encryption tools, such as PGP. The payment service debits the relevant sum from the credit 25 card account, and generates a number of payment tokens, say 1000 tokens of value 1p. These are encrypted using the public key algorithm and returned to the user via the internet connection, together with a key which is unique to the user. Each token comprises, in this example, a 64 bit random hexadecimal number, drawn from a large list of n random numbers $R = (r_0, r_1, r_2, \dots, r_{n-2}, r_{n-1})$ at the 30 payment service. For each user, the payment service keeps two pieces of secret information k and s . k is a random key for use with a symmetric block cipher. s is a random security parameter, where $(0 \leq s \leq n-1)$ taken at random from the range

(0..n). There is also an integer index variable i . Its secrecy is not essential although its integrity is important.

The carnet which is stored at the client terminal holds a list of m random numbers $T = (t_1, t_2, t_3, \dots, t_{m-1}, t_m)$, where $(m \ll n)$. Table 1 below shows part of the list of random numbers held in the carnet. The payment service derives T from R when the user registers. T is an encrypted subset of R , such that;

$$t_x = E(k, r_{(s+x \bmod n)})$$

where $E(a, b)$ denotes encryption of b using key a . The variable i is initialised to 0 at this point. The carnet also contains a copy of this variable.

To make a single micropayment, the user sends t_i (the i th payment token) to the payment service. The payment service checks to ensure that

$$D(k, t_i) = r_{s+i \bmod n}$$

and the payment succeeds. The carnet and payment service both increment their copies of i .

Subsequently, the user establishes an Internet connection with the merchant. This may be, for example, in order to retrieve HTML pages of news items which have been retrieved by the Merchant using a search engine and search criteria which were previously specified by the user. In response to a request to download the HTML pages, the Merchant requests payment of, e.g., 10p. Typically, the returning of the request for payment will be automated using CGI scripts running on the HTML server at the Merchant.

In response to the request from the Merchant, the user issues the tenth token from the Carnet. Since the tokens are in a determined sequence it is sufficient to transfer, e.g. just the tenth token in lieu of all ten tokens required to make the payment total. This token is modified by the carnet module by an XOR operation with a PIN which the user is required to input at the terminal in order to authorise the issuing of tokens. The modified token is then transmitted on the internet connection to the Merchant.

The Merchant has an open internet connection to the payment service. When the tokens are received from the user, the Merchant transmits to the Payment service a tuple containing the payment tokens, together with a corresponding set of authentication tokens. The Merchant, also termed herein the "third party service provider" (TPSP), is issued with a set of authentication tokens,

which are essentially the same as the payment tokens issued to a user, but which function like one-time passwords. The set of authentication tokens A is an encrypted subset of R , such that

$$a_x = E(k, r_{s+x \bmod n})$$

- 5 where k is a secret encryption key, s is a random security parameter and n is the size of the random database R .

The TPSP collects payment tokens t_i from the user as previously described. To verify each token, the merchant forwards the token together with a_j (the next authentication token) to the payment service. The payment service verifies that
 10 both t_i and a_j are the expected values before returning confirmation. By counting the (t_i, a_j) pairs submitted successfully by the merchant, the payment system has a record of the amount that merchant is owed.

The payment service returns an authentication token in the authentication message. If the payment token was submitted by the merchant with
 15 authentication token a_j , then the payment server returns $a(j+1)$, that is the next authentication token in the sequence held by the merchant, to indicate that authentication was successful. The payment server may transmit the complement of $a(j+1)$ to indicate that authentication has failed.

The messages exchanged by the user, the merchant and the payment
 20 server are all short and are suitable for inclusion, for example, in HTTP (hypertext transfer protocol) or MIME (multipurpose internet mail extensions) headers.

The payment system as described so far is potentially open to several types of failure. Mechanisms to handle these are included in this implementation. There are two main classes of recoverable failure:

- 25 1) A carnet or merchant module's index variable may get out of step with the payment service's. These errors are handled by a resynchronisation scheme. If a received token does not match its expected value, the payment service reads forward through the sequence of random numbers looking for a value that does match. A read forward limit is set as a system parameter. If a matching token is
 30 found within the range set, the payment system accepts the token. It records the number of tokens skipped for each carnet. A carnet can recover the tokens later (e.g. when the number of tokens remaining gets low) by performing a refresh.

Optionally, this scheme may be elaborated by using a sliding window which defines the range of values against which a token may be elaborated.

- 2) The list of tokens or the index variable in either a carnet could become lost or
5 damaged. A carnet can be refreshed, making its state consistent with the
information held in the payment service. When a carnet is reinitialised, a new set
of payment tokens is constructed for the user. The number of tokens is equal to
the number of unspent tokens in the previous carnet, plus the number of tokens
lost through synchronisation
10 errors.

There is potentially a third type of failure which is not recoverable. Corruption of
the payment service's information would be fatal and must be prevented. This is
accomplished by building redundancy into the payment service, for example by
15 mirroring data from one payment server at another server at a different site, and by
using robust storage technologies such as RAID.

A further feature of the present implementation is that each carnet issued
to a user is provided with a digital signature. The signature is formed in a
conventional fashion using a public-key encryption algorithm. It is particularly
20 preferred that DSA or another El Gammal variant should be used for this purpose.
The digital signature then may be used for dispute resolution. For example, if the
user purchases tokens which, as a result of a system failure, can not be spent,
then the user may claim a refund from the party who issued the carnet. That party
can ensure that the carnet is genuine by checking that the digital signature is that
25 generated when the carnet was issued.

It will be understood that the above application of the QuickPay
technology is described by way of example only, and that the invention may be
used in a wide variety of different contexts, wherever a secure and efficient
payment service is required. Although the invention has particular advantages in
30 the context of internet services, because of the low messaging costs, it may also
be used in other areas. For example, the carnet and associated client applications
might be stored on a smartcard including a microprocessor and non-volatile
memory. The merchant platform in this case might be accessed via a card reader

at an EPOS (electronic point of sale) terminal. Figure 9 shows such a system, in which the carnet on a smartcard 901 is in a card reader 902 connected to a merchant platform 903. The merchant platform is connected via the PSTN to the payment server 904. The card might also be used to purchase telephony services
5 via a card reader attached, for example, to a public telephone. In this case, the existing telephony network signalling channels, such as NUP (national user part) and ISUP (ISDN user part) in the UK PSTN, may be used for the connections between the different platforms, in place of TCP/IP signalling in the first example described above.

10 In a further example, QuickPay is used to pay for use of BT's NetSumm service. NetSumm (Trademark) is a networked text summarisation tool. With it, internet users can produce summaries of English language World Wide Web (WWW) documents. The service is accessed via a Web browser.

The user first visits the QuickPay web site. First they set up an account by
15 completing a form supplying a username and a password. The username should be a valid email address, the password can be any character string up to 64 bytes long. The user then visits a page where they can purchase a QuickPay carnet . The user completes another form on which they specify the type of platform they use and the size of the carnet they require. When they click on the 'Buy' button,
20 the payment service creates a new QuickPay carnet and downloads it to the client. The user then installs their new carnet . They can specify the location that the QuickPay client will be installed at and they must supply a PIN to protect the payment tokens. From the NetSumm homepage the user specifies the URL (uniform resource locator, i.e. the address) of a WWW document to be
25 summarised. NetSumm returns a digested version of the document. All hyperlinks in the document are modified so that clicking them returns a summary of the target document. NetSumm also presents additional controls to the user to allow them to change the summarisation rules.

Use of NetSumm is charged at £0.002 (0.2 pence) per summary. Every
30 time a page is submitted, NetSumm charges the user two QuickPay tokens before returning the result (the token value in this example is 0.1p). A single call to the QuickPay merchant module allows the NetSumm application to request these tokens and get a result indicating if the payment was successful or not. Before

using NetSumm, the user must run their QuickPay client. This is an application separate from the user's web browser. It has a single small window which gives a visual indication of the number of tokens remaining. The window also contains some controls to allow the user to quit the application, change the PIN, update the
5 carnet etc. If the user fails to start the client before using NetSumm, they will receive an error page from the service when they attempt to summarise a text. The first time the NetSumm service attempts to collect tokens from the QuickPay client, the user is presented with a dialogue box asking them to confirm the payment. They must enter their PIN to unlock the payment tokens and they can
10 choose how to handle any subsequent payment requests from this merchant. If the user chooses to allow subsequent payments, NetSumm will be able to collect additional tokens without user intervention. Alternatively, the user may choose to accept or reject each request individually. These rules persist until the user closes the QuickPay client application.

15 When the user has exhausted their supply of payment tokens, they obtain more by visiting the QuickPay web site, where they follow a 'Refill Carnet' link . Refilling the carnet is very similar to purchasing the first one, except that any unused tokens from the previous carnet are added to the user's new one. The user must enter their authentication details (username and password) to purchase
20 additional tokens. They must also provide payment details (i.e. a credit card number) and select the type and size of carnet they want. Any unused tokens from the previous carnet are added to the new one before it is delivered and the user's old carnet becomes invalid.

An implementation of the invention will now be described in further detail.

25 Figure 1 shows the main software components in this implementation of QuickPay. The payment service and merchant module are back end applications, with little or no user interaction. These components are, in this example, implemented as UNIX client and server applications and should run on most UNIX platforms. The carnet does interact with its user and may be implemented in a
30 variety of forms on several types of computer. A generic UNIX client is described here, which may be run on UNIX workstations, or may be used as the basis of stand alone clients ported to other platforms. As a further alternative, the

QuickPay client may be integrated with other applications, for example WWW and email clients.

Payment Service

The payment service comprises the clearer, the creator, the updater and four
5 databases.

Databases

All four databases are simple collections of records. In the URD (User Registration Database), entries are indexed by an alphanumeric string (the username). In all other databases the entries are indexed sequentially by number. There is no
10 requirement for the allocation of indices to be contiguous. For example, carnet IDs are allocated at random from the range [0 .. 264-1] and only a fraction are likely to be allocated at any one time.

User Registration Database

```
15      struct {
          char passwd[];          /* Password */
          verylong cid;          /* Carnet ID */
      }
```

The URD holds information about QuickPay users. Its main use is for authentication of users who are performing administrative functions.

20 Entries in the URD contain a password and a carnet ID.

The Carnet Information Database

```
      struct {
          verylong key;
          verylong dbid;
25      u_long start;
          u_long size;
          u_long index;
          u_long resyncs;
      }
```

30 The CID (Carnet Information Database) holds state information for the user carnets and for the merchant modules. This database is used by the clearer to identify the next expected payment or authentication token. It also allows the updater to reconstruct any set of payment or authentication tokens from the RND.

The key is the value used in the keyed hash to construct the tokens. The database
35 ID, start offset and size fields identify the source of the random data in the RND. The token index is the index of the next expected token and lies in the range [0 ..

size]. The resync (resynchronisation) counter is a record of number of tokens that have been skipped in order to clear a received token.

Merchant Information Database

```

5      struct {
          u_long accepted;
      }

```

The MID (Merchant Information Database) holds the record of transactions cleared for each merchant. Where an entry for a carnet in the CID is for a merchant, there
 10 is a corresponding entry in the MID.

The Random Number Database

```

      struct {
          u_long size;
          verylong tokens[];
15      }

```

The RND (Random Number Database) holds a pool of random data which is used to construct sets of payment and authentication tokens. Each entry in the database contains a random sequence of 64 bit values. The size field indicates how big this sequence is. The RND is populated with data from a hardware random number
 20 generator, as follows: The program RandomCommServer can be divided into three sections. The random generator comes under the source file name Random Generator.cpp with header file Random Generator.h the random tester uses source files Main Tester.cpp, tester.cpp, and Random IO.cpp and uses header files Main Tester.h tester.h, stats.h and Random IO.h. The remaining parts serve as the
 25 windows interface and the server. The main source files are RCTServerDoc.cpp, ServerSocket.cpp and "CommSocket.cpp".

1. Generating the data.

The random numbers are generated using the a hardware random number generator manufactured by BT and known as Lektor 3900, in blocks of 256K
 30 (=2,097,152 bits). The random number generator is a card which fits into a personal computer, and which contains a diode and an amplifier which amplifies shot noise on the diode to create a random stream of binary values. The generator has a slight bias, so to correct this, the data is compressed by changing two bits of data into one bit or no bits. 00s and 11s are discarded, 01 becomes 0 and 10

becomes 1. This means that three-quarters of the data is lost, but the bias is removed. This data is then saved on disk under the name "raw.dat".

2. Testing the data

The following tests are used on the blocks of 256K

- 5 (i) Bits Test: tests the number of zeroes against the number of ones.
- (ii) Serial Test: tests the frequencies of 00, 01, 10, and 11.
- (iii) Runs Test: tests the occurrences of long sequences of 0..0 or 1..1.
- (iv) Poker Tests of size 2 to 8: tests the occurrences of 2- to 8-bit combinations.
- (v) Autocorrelation Test for periods 1 to 1000: tests for periodicity in the data.
- 10 (vi) Maurer Universal Statistical Test: looks at the potential for the compression of data.

In addition, the bits test is calculated cumulatively for all the generated data. This checks for an overall bias, which may not be evident in tests of data of only 128K.

15 3. Acceptance/ Rejection of the data.

- The data is tested at the 5% and 1% significance levels. This should mean that the number of tests failed at the 5% level should be 1 in 20, and the number at the 1% significance level should be 1 in 100. If the data fails at significantly more of these 1011 tests (80+ at the 5% level, 15+ at the 1%, or 1 "off-the-scale" error) then the data will be rejected. In addition a score of the failures at a particular test will be kept. If one test persistently fails, then this will be flagged.
- 20

4. Storage of Data and Use as Server

- Once the data has been accepted as genuinely random, the 256K file is split into eight 32K files and stored under the filename "use_n.dat" (where n is a three-digit number). This will be available to any client machine that connects to the server.
- 25
- The server sends out files, using the oldest first. Once a file has been sent to a client, it is deleted from storage. At any one time, up to 1000 32K files will be available to clients. If it is feared that the random data can be accessed while it is being sent, then it is recommended that the client encrypt the received random
- 30 data.

Creator

The creator is invoked to create a new carnet for a user. Figure 2 gives a flow diagram for this process. It assumes that if this is the user's first carnet, an account has already been created in the URD.

The creator picks a new carnet ID, key, database ID and start offset at random and
5 adds these details to the CID. It then reads tokens from the RND, encrypts them using the key and packages them (together with the client software) in an installer for distribution.

Clearer

Figure 3 gives a flow diagram for the main execution loop of the payment service
10 clearer module. Its function is to process payment tokens collected by QuickPay merchants. The clearer listens for requests from merchants on the network. When a request is made, the clearer reads the tuple {im, tm, ic, tc, n }, (where im is the identity of a carnet at the merchant, tm is a token from that carnet at the merchant, ic is the identity of a carnet at the client, tc is a token from that carnet
15 at the client and n is the number of tokens from the client) and checks that it is properly formatted. The clearer ignores improper requests. The clearer checks {im, tm} and {ic, tc, n} (see below). If both of the checks are passed, the payment service increments the accounting record for the merchant to indicate the acceptance of another n tokens. An authentication message is returned to the
20 merchant. Optionally, the authentication message may take the form of the next in a sequence of authentication tokens previously issued to the merchant. The complement of the token value may be transmitted when authentication has failed. Figure 4 describes the process of verifying the {i, t} or {i, t, n} tuple. The clearer starts by reading the key k, database ID d, start offset s and token counter c from
25 the CID. Using s, c and n, the clearer determines the position of the next expected token for the carnet within the raw database d. It reads this token and forms the keyed hash using k. If the result is the same as the supplied value, the token is accepted.

Sometimes, the condition will arise where the index pointer in the carnet gets
30 ahead of that in the payment service. This happens as a result of tokens being collected but not cleared. The clearer has a mechanism for dealing with this problem. If the first token retrieved from the RND does not match the clearer searches forward, trying the next token, then the next and so on. The clearer will

accept a token which matches within 100 attempts (this figure is variable according to the needs of a particular implementation). For every token in the RND which is skipped, the clearer increments the resync counter for the carnet in the CID. The clearer reports the success or failure by returning a numeric code. As
5 noted above, the basic scheme may be made more flexible by the use of a sliding window rather than simply a fixed range forward from the current index pointer at the payment service.

Figure 10 illustrates the sliding window scheme. The figure shows how the $\{i, t\}$ and $\{i, t, n\}$ tuples are verified. The clearer starts by reading the key k ,
10 database ID d , start offset s and token counter c for the carnet i , from the CID. The values s , c and n are used by the clearer to determine the position of the next expected token for the carnet within the raw database d . It reads this token and forms the keyed-hash using k . If the result is the same as the supplied value t , the token is accepted. To allow for some slight re-ordering of transactions, the clearer
15 maintains a transaction window surrounding the position of the next expected token. Instead of simply testing the value at this position, the clearer will accept a token if it occurs anywhere within the window.

Initially, the window begins at the position of the next expected token. If transactions arrive out of sequence, the clearer searches forward within the
20 window in an attempt to clear the token. If a match is found the clearer records the fact that a particular token has been cleared to prevent double spending. Positions in the stored sequence which correspond to cleared tokens are shown shaded in the Figure. The size of the window is a configurable system parameter. The clearer must store a flag to indicate the state of every token within the
25 window, so the size of the window influences the storage requirements of the clearer.

As tokens are spent, the window moves forward along the list of numbers in the raw database. The window moves under two situations;

1. As tokens are cleared at the back of the window.
- 30 2. When the clearer must search forward beyond the front of the window to clear a token.

In the case of (2), tokens are lost at the back of the window as the window moves forward. The clearer additionally maintains a count of tokens lost in this way (the resync count).

Optionally, an implementation of the QuickPay system may opt not to use windowing by configuring a window size of 1. In this way the clearer avoids maintaining state flags for items in the window, but is unable to clear transactions out of sequence.

Updater

- 10 The updater is similar to the creator, except that it rebuilds carnets from existing entries in the URD and CID. It is used to restore a client or merchant's token database should it become damaged or out of sync with the payment service. The updater can also be used to recover tokens dropped (never cleared) by the payment service. Figure 5 gives a flow diagram for the updater.
- 15 The updater reads the carnet size *s*, token counter *c* and resync counter *r* from the CID. It calculates the number of tokens remaining and generates a new key, database ID and start index. The rest of the algorithm is as for the creator.

Merchant Module

- 20 The merchant module comprises the collector, administration functions and the authentication token database.

The Collector

- The collector is invoked when the merchant wishes to collect a payment from a user's carnet. Figure 6 gives a flow diagram for the execution of the collector. In
- 25 step 61 the collector module establishes a connection to the wallet in the Payer module (described below). In step 62 the Nth token is requested from the Payer, where N is determined by the number of tokens required to give the necessary monetary value for the transaction. In step 63 the response from the Payer is received. The response is tested in step 64, and if the request was refused then
 - 30 failure is signalled in step 610. Otherwise, in step 65 the merchant connects to the payment service. In step 66 the tuple containing the following values is transmitted: Merchant ID (MID), Authorisation token (AT), N, Wallet ID (WID), Payment Token (PT). In step 67 a response from the payment server is received

and in step 68 this response is tested. If the request for authorisation is not refused, then in step 69 success is signalled to the Payer. Otherwise, if the request for authorisation was refused, then in step 610 failure is signalled.

Administration Functions

- 5 The merchant module includes administration functions. These maintain a count of how many unused authentication tokens remain, and send a request for further tokens to the payment service when that number falls below a predetermined threshold.

Carnet

- 10 The carnet comprises the payer, the payment token database and the administration functions. There is also an installer responsible for initial installation on the client's platform.

The Payer

- 15 The payer is responsible for handling requests for payment tokens. It receives requests from the network and decides whether to honour them. This decision is based on a number of factors, including user input, the number of tokens remaining and the history of previous transactions. Figure 7 gives a flow diagram for the payer.

Administration Functions

- 20 As for the merchant module, the carnet includes a mechanism for refreshing the carnet and adding new tokens. In addition, the payment tokens stored on the user's machine are protected by a PIN (personal identification number).

The Installer

- 25 The installer sets up the QuickPay client on the user's machine. It creates a number of files and directories and allow the user to PIN-protect the newly installed payment tokens. The QuickPay client in this example is distributed as a compressed tar file for UNIX systems

TABLE 1

Wallet ID	0xbfb0a569fcb61f0
0	0x502a1d6c3351d07c
1	0xc30ce637bd6091b3
2	0x6420dda77f68a7db
3	0x85e354f3229f48f1
4	0x81399aa9cb5a87ae
5	0x26c5724d56794a57
6	0xccf8b8add7765a7f
7	0x4ab802396f89d6ed
8	0x5f9a820ce1e667e7
9	0x1e5a225186e01181
10	0xa32cd70d233c50d7
11	0x956e712aaa29cf4d
12	0xb92d9eecd3f16ec
13	0x11493c8bfcee102e
14	0xa870480af34d6778
15	0x9cc5f5945e153bba
16	0xb6abd8c968c47312
17	0x7cfb9ab0555c58c5
18	0x1051e7a417766c01
19	0x5d5e95685bc01249
...	

CLAIMS

1. A method of operating a digital payment transaction system comprising:
 - a) storing at a payment server a sequence of random numbers;
 - 5 b) issuing to a user a set of digital payment tokens comprising a sequence of digitally encoded random numbers derived from the said stored sequence of random numbers;
 - c) transferring a payment token from the user to a merchant platform;
 - d) transferring from the merchant platform to the payment server the
10 payment token received from the user in step (c);
 - e) at the payment server, authenticating the token by comparing the value of the random number of the token from the merchant platform and a value derived from a corresponding position in the stored sequence of random numbers;
and
 - 15 f) subsequently communicating an authentication message from the payment server to the merchant platform.
2. A method according to claim 1, in which the payment server is remote from the merchant platform and the merchant platform communicates over a
20 communications network with payment server.
3. A method according to claim 1 or 2, in which the set of digital payment tokens is derived from the sequence of random numbers stored at the payment server by selecting part of the said sequence and encoding part of the said
25 sequence with a key which is specific to the user.
4. A method according to claim 3, in which the part of the said sequence is encoded by a symmetric block cipher.
- 30 5. A method according to any one of the preceding claims, in which the user is remote from the merchant platform and in step (c) transfers the payment token to the merchant platform via a communications network.

6. A method according to any one of the preceding claims, in which the step of authenticating the digital payment token comprises:

i) attempting to authenticate the digital payment token against a value at a position in the sequence of random numbers stored at the payment server; and

5 ii) when the token is not authenticated in step i), attempting to authenticate the digital payment token against one or more other values in the stored sequence, which other values fall within a predetermined maximum distance from the said position;

and in which in step (f) the authentication message indicates that the
10 authentication is successful when the token is successfully authenticated in either step (i) or step (ii) .

7. A method according to any one of the preceding claims, in which in step (d) the merchant communicates together with the payment token an authentication
15 token from a sequence of authentication tokens issued by the payment server.

8. A method according to claim 7, in which the payment server automatically updates a merchant account record after authenticating a validated payment token and authentication token received from the merchant platform.

20

9. A method according to any one of the preceding claims further comprising:

f) maintaining at the payment server a record of the current state of the set of digital payment tokens; and

g) when the digital payment tokens issued to the user are lost or
25 corrupted, communicating data from the payment server to the user and thereby updating the set of digital payment tokens to a state corresponding to that recorded at the payment server.

10. A method according to any one of the preceding claims further comprising:
30 issuing the user an identification number (PIN);

modifying, using the identification number, the numbers derived from the said stored sequence of random numbers ;

and further modifying, using the identification number, the digital payment token which is transferred to the merchant in step (c).

11. A method according to claim 10, in which, in the steps of modifying and
5 further modifying the digital payment token, the digitally encoded value derived from the stored sequence of random numbers is combined with the identification number using a Boolean logic operation.

12. A method according to claim 11, in which the said Boolean logic
10 operation is XOR, and, when the same identification number is used in the steps of modifying and future modifying, the step of further modifying the digitally encoded value recreates the original value derived from the said stored sequence of random numbers.

13. A method of operating a digital payment transaction system comprising:
15 a) issuing to a user a set of digital payment tokens comprising a sequence of digitally encoded random numbers
b) issuing to a merchant a set of authentication tokens comprising a sequence of digitally encoded of random numbers;
20 c) transferring a digital payment token from the user to a merchant;
d) transferring the said digital payment token, and with the digital payment token transferring an authentication token, from the merchant to a payment server;
e) authenticating the digital payment token and the authentication
25 token against records stored at the payment server; and
f) returning an authentication message to the merchant.

14. A method according to claim 13, in which the step of authenticating the digital payment token comprises:
30 i) attempting to authenticate the digital payment token against a value at a position in a sequence of random numbers stored at the payment server; and
ii) when the token is not authenticated in step i), attempting to authenticate the digital payment token against one or more other values in the

stored sequence, which other values fall within a predetermined maximum distance from the said position;

and in which in step (f) the authentication message indicates that the authentication is successful when the token is successfully authenticated in either
5 step (i) or step (ii) .

15. A method according to any one of the preceding claims, in which an authentication token is returned by the payment server with the authentication message.

10

16. A method according to claim 15, in which the said authentication token corresponds to one of a plurality of authentication tokens previously issued to the merchant.

15

17. A merchant platform for use in method according to any one of the preceding claims, the merchant platform comprising:

means for receiving from a user a digital payment token which comprises a digitally encoded random number

20

means for transferring the digital payment token to a payment server; and

means responsive to an authentication signal issued by the payment server when the said digital payment token is successfully authenticated.

18. A merchant platform according to claim 17, further comprising a store
25 programmed with a sequence of authentication tokens issued by the payment server, in use an authentication token being returned to the payment server with each digital payment token.

19. A payment server for use in a method according to any one of claims 1 to
30 16, the payment server comprising

a data store programmed with a sequence of random numbers

means for issuing digital payment tokens which comprise a sequence of digitally encoded random numbers derived from the sequence in the data store

means for receiving digital payment tokens returned by a merchant platform;

means for authenticating the returned digital payment tokens; and

means for outputting an authentication message to the merchant platform.

5

20. A client platform for use in a method according to any one of claims 1 to 16, the client platform comprising;

means for receiving from a payment server which is remote from the client platform digital payment tokens which comprise a sequence of digitally encoded

10 random numbers;

a token store for storing the said digital payment tokens; and

payment means for issuing a token from the sequence to a merchant platform.

15 21. A communications network including one or more of a merchant platform according to claim 17 or 18, a payment server according to claim 19, and a client platform according to claim 20.

Fig.1.

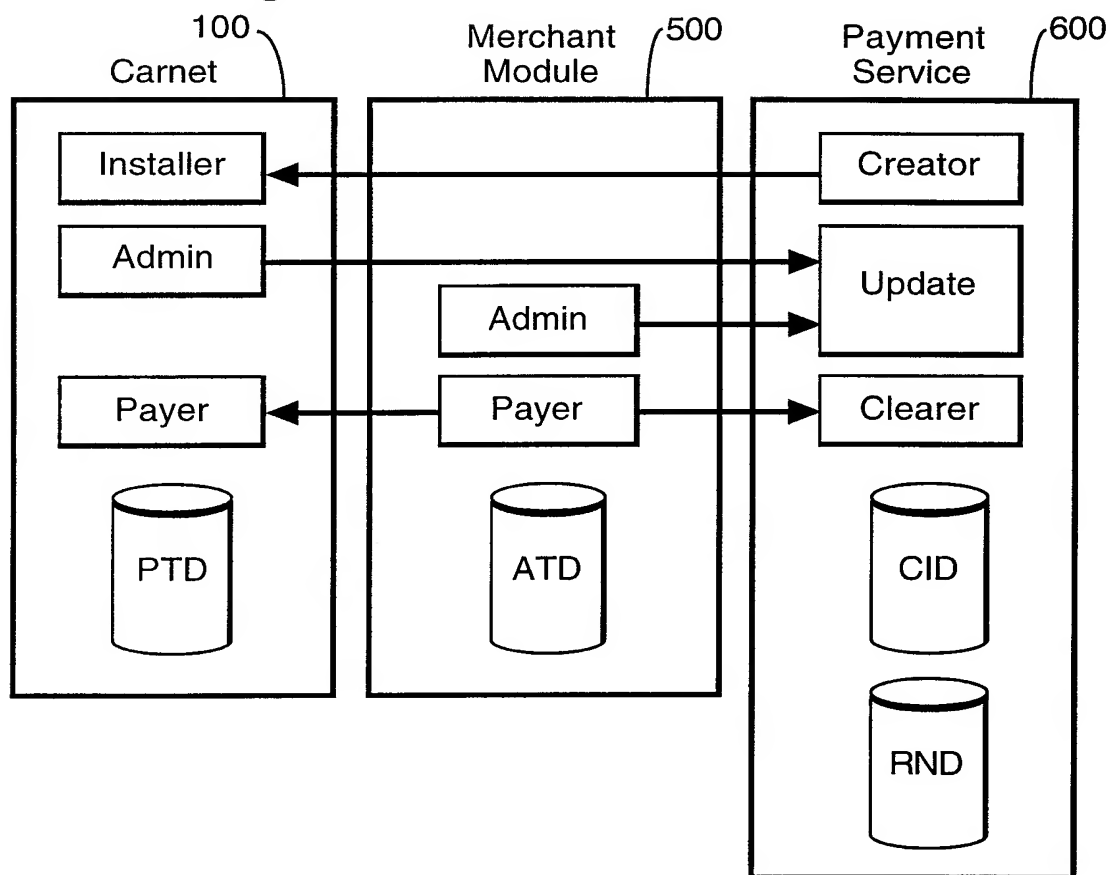


Fig.2.

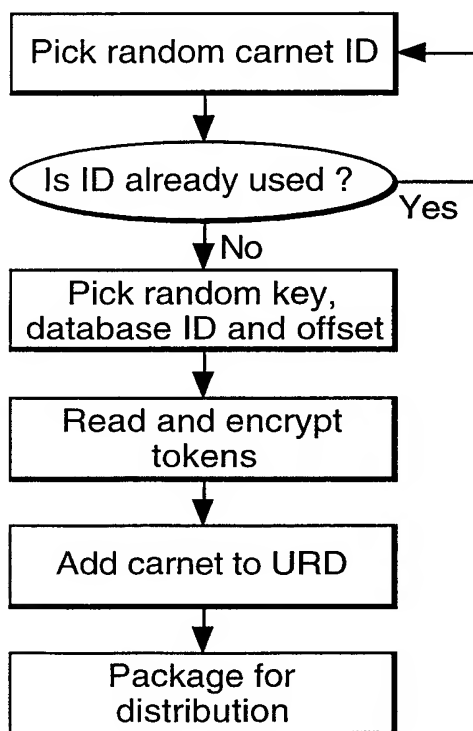


Fig.3.

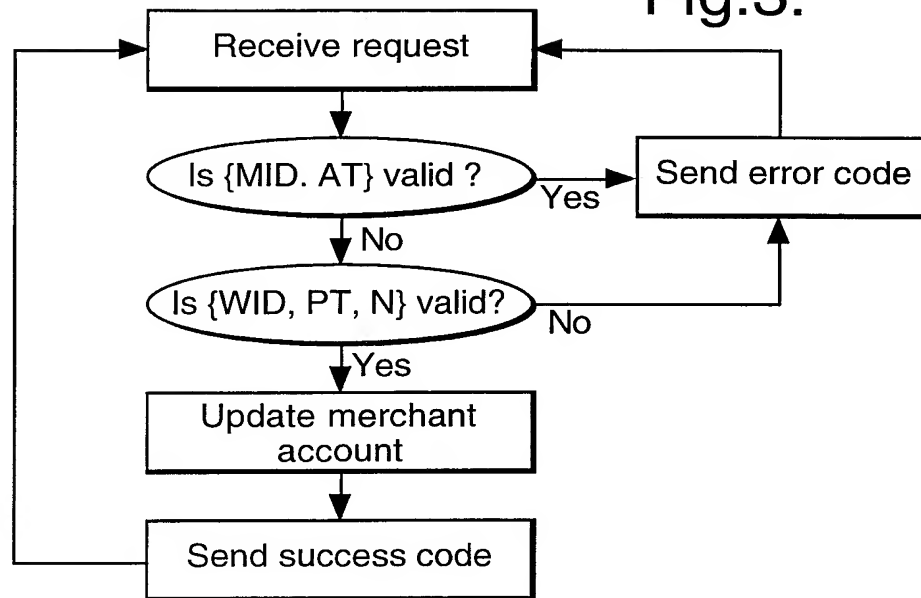


Fig.4.

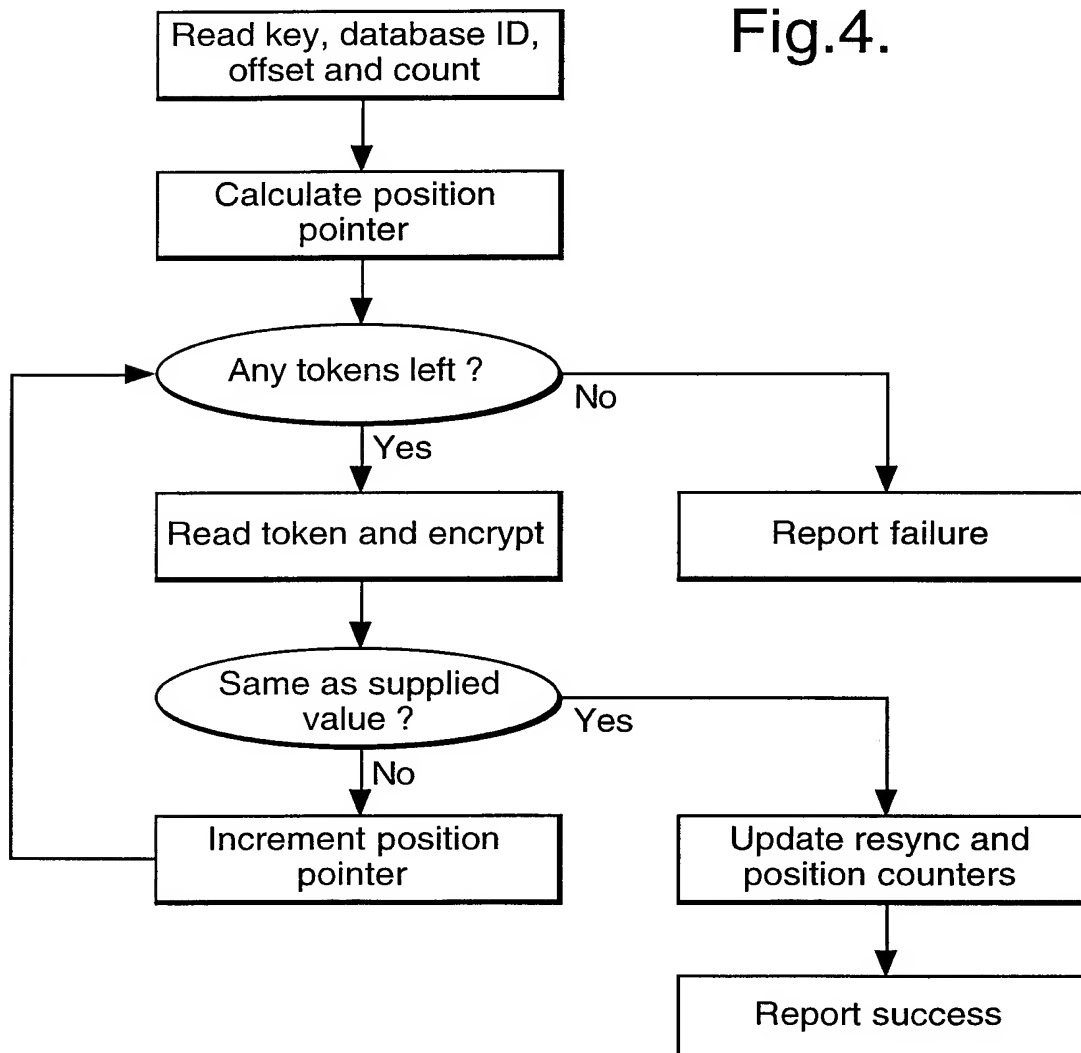


Fig.5.

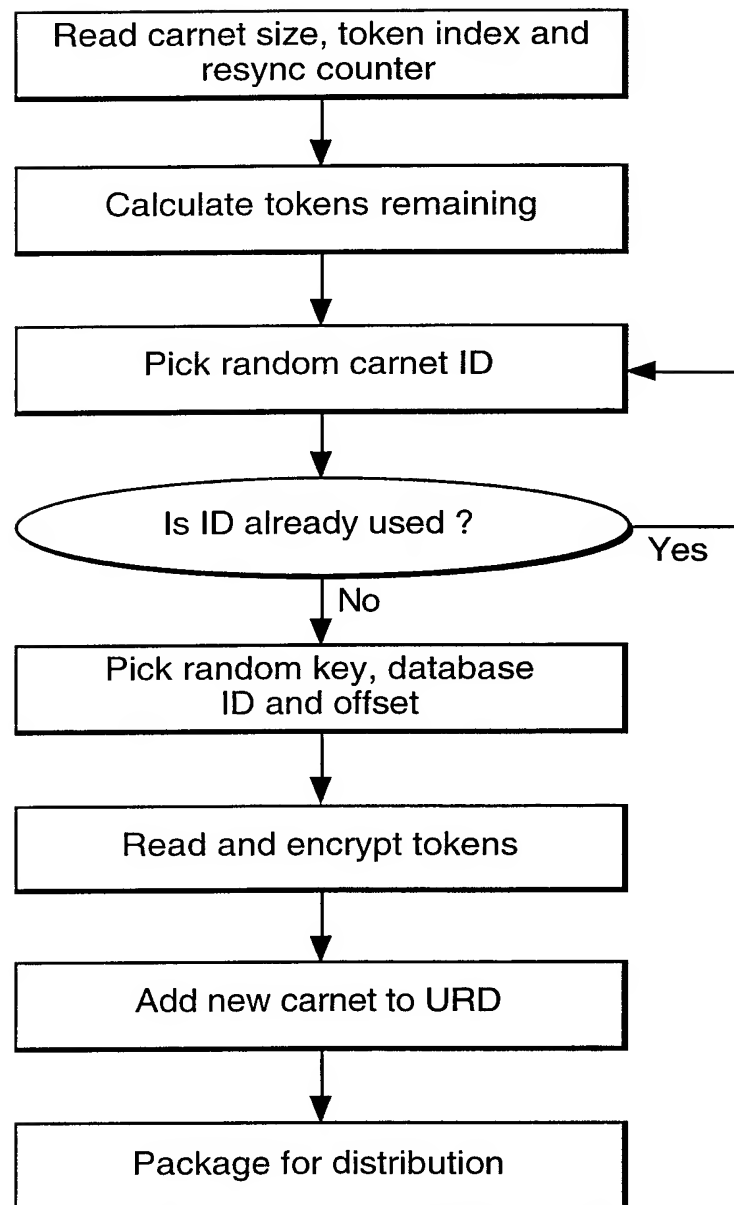
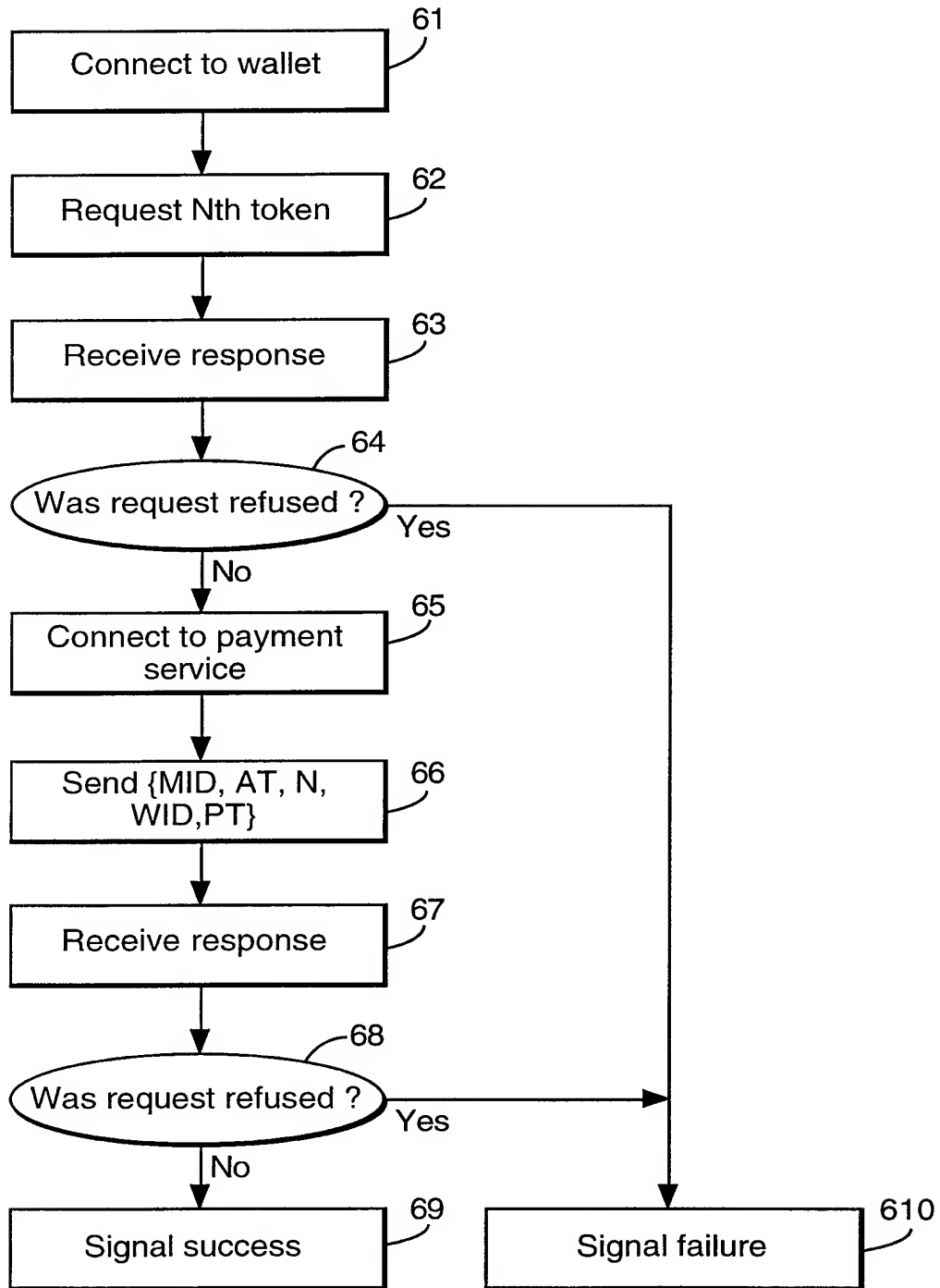


Fig.6.



5/6

Fig.7.

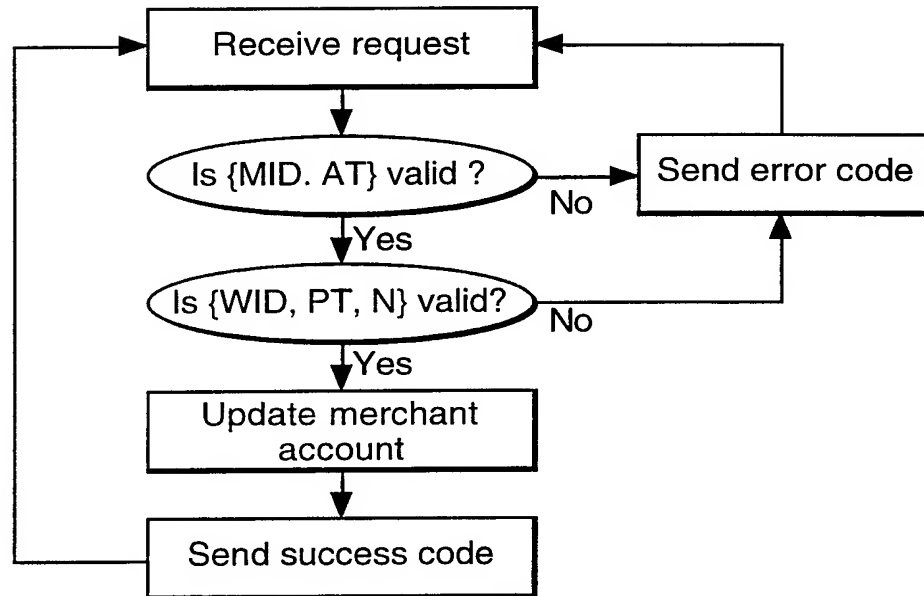


Fig.8.

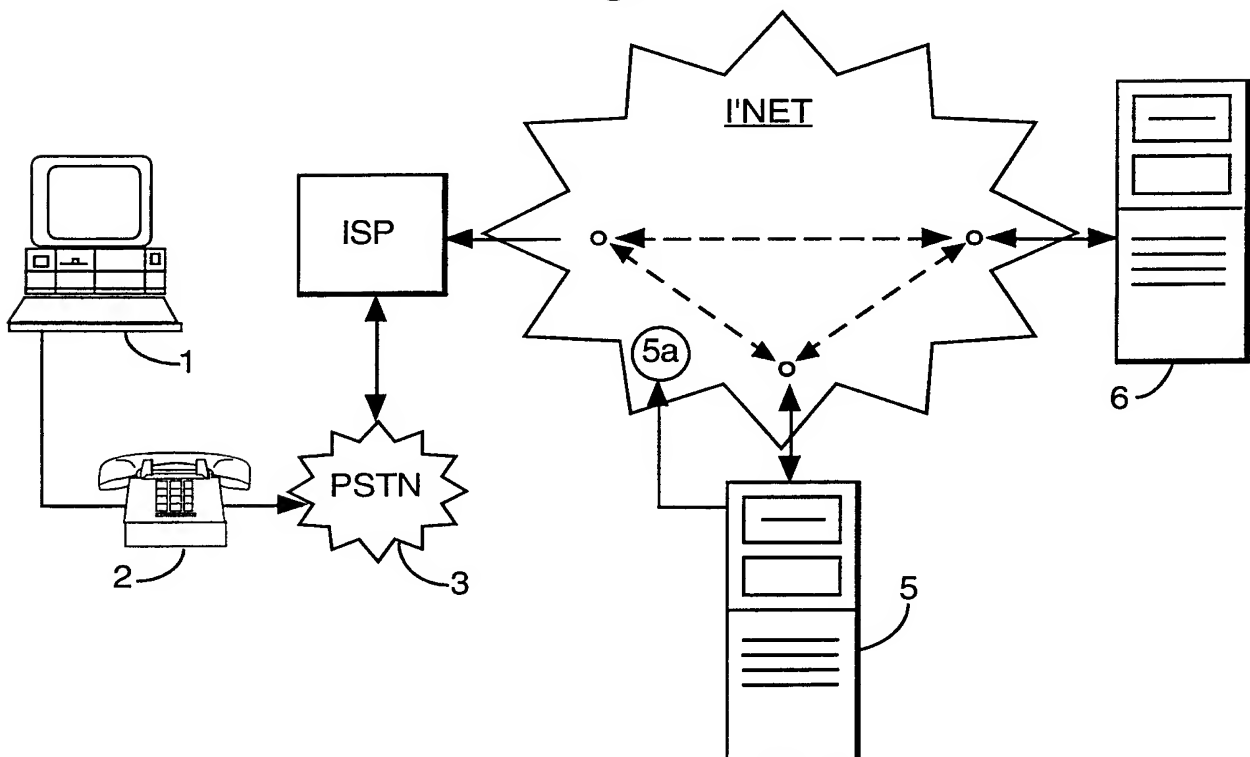


Fig.9.

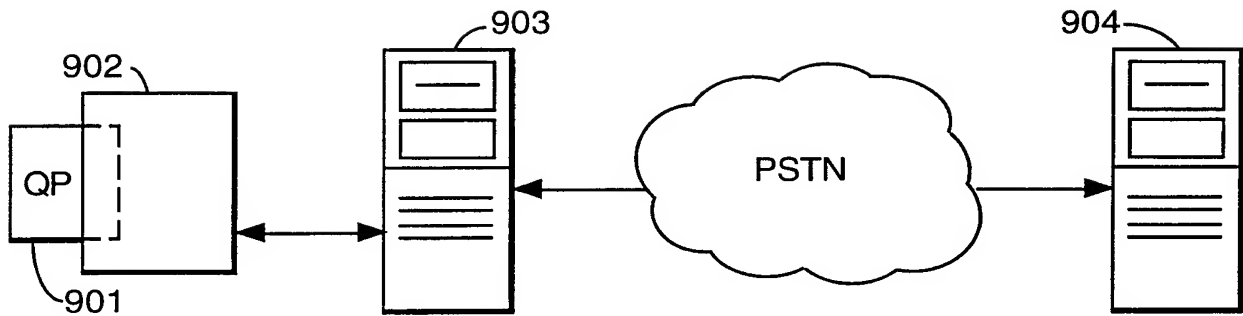
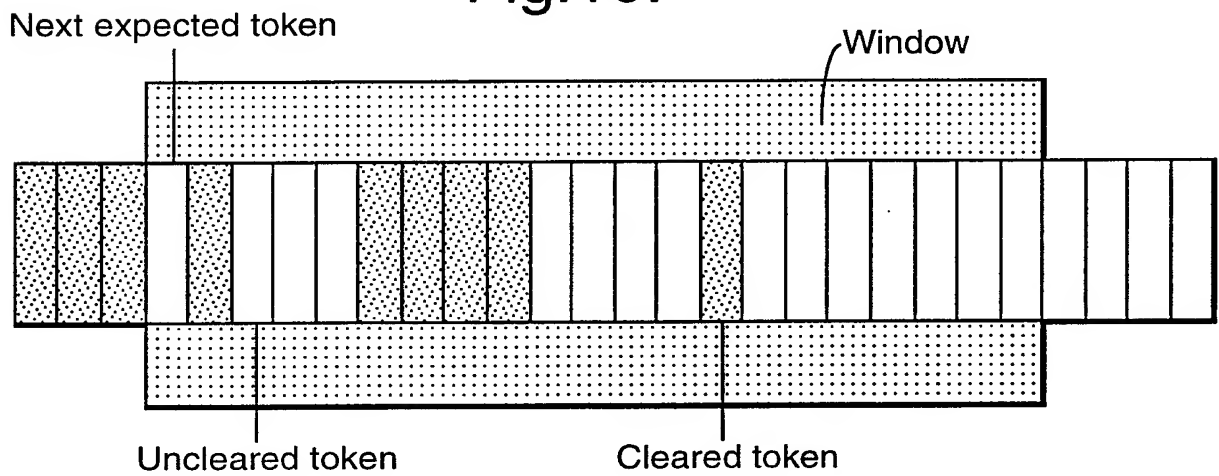


Fig.10.



INTERNATIONAL SEARCH REPORT

International Application No

PCT/GB 97/03116

A. CLASSIFICATION OF SUBJECT MATTER

IPC 6 G07F7/08 G07F7/10

According to International Patent Classification(IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)

IPC 6 G07F

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category *	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
X	EP 0 507 669 A (FRANCE TELECOM ;TELEDIFFUSION FSE (FR); PHILIPS NV (NL)) 7 October 1992	1,11,12
A	see claim 1; figure 1	2-10, 13-16
P,X	WO 97 03423 A (DIGITAL EQUIPMENT CORP) 30 January 1997 see claim 1; figure 11	1,11,12
P,X	US 5 590 197 A (CHEN JAMES F ET AL) 31 December 1996 see claim 1; figure 1	1,11,12
P,A	WO 97 03410 A (EGENDORF ANDREW) 30 January 1997 see claim 1; figure 1	1-16
	-/--	



Further documents are listed in the continuation of box C.



Patent family members are listed in annex.

* Special categories of cited documents :

"A" document defining the general state of the art which is not considered to be of particular relevance

"E" earlier document but published on or after the international filing date

"L" document which may throw doubts on priority claim(s) or which is cited to establish the publication date of another citation or other special reason (as specified)

"O" document referring to an oral disclosure, use, exhibition or other means

"P" document published prior to the international filing date but later than the priority date claimed

"T" later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention

"X" document of particular relevance; the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone

"Y" document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art.

"&" document member of the same patent family

Date of the actual completion of the international search

13 February 1998

Date of mailing of the international search report

24/02/1998

Name and mailing address of the ISA

European Patent Office, P.B. 5818 Patentlaan 2
NL - 2280 HV Rijswijk
Tel. (+31-70) 340-2040, Tx. 31 651 epo nl,
Fax: (+31-70) 340-3016

Authorized officer

Kirsten, K

INTERNATIONAL SEARCH REPORT

International Application No

PCT/GB 97/03116

C.(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT

Category	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
A	WO 96 08783 A (FIRST VIRTUAL HOLDINGS INC) 21 March 1996 see claim 1; figure 1 ----	1-16
A	GB 2 241 414 A (ENFRANCHISE SIXTY LTD) 28 August 1991 see claim 1; figure 1 ----	1-16
A	US 5 267 314 A (STAMBLER LEON) 30 November 1993 see claim 1; figure 16 -----	1-16

INTERNATIONAL SEARCH REPORT

Information on patent family members

International Application No

PCT/GB 97/03116

Patent document cited in search report	Publication date	Patent family member(s)	Publication date
EP 0507669 A	07-10-92	FR 2674976 A DE 69211407 D DE 69211407 T JP 6325055 A US 5305383 A	09-10-92 18-07-96 06-02-97 25-11-94 19-04-94
WO 9703423 A	30-01-97	EP 0796480 A JP 9510814 T	24-09-97 28-10-97
US 5590197 A	31-12-96	NONE	
WO 9703410 A	30-01-97	AU 5986596 A	10-02-97
WO 9608783 A	21-03-96	AU 3630995 A CA 2199942 A EP 0791202 A	29-03-96 21-03-96 27-08-97
GB 2241414 A	28-08-91	AU 7238991 A WO 9112680 A	03-09-91 22-08-91
US 5267314 A	30-11-93	US 5524073 A US 5646998 A US 5555303 A	04-06-96 08-07-97 10-09-96